

# TANGO: Secure Collaborative Route Control across the Public Internet

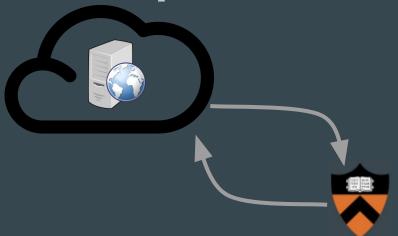
Henry Birge-Lee, Sophia Yoo, Benjamin Herber, Jennifer Rexford, Maria Apostolaki



Is today's Internet good enough for performance-critical applications?

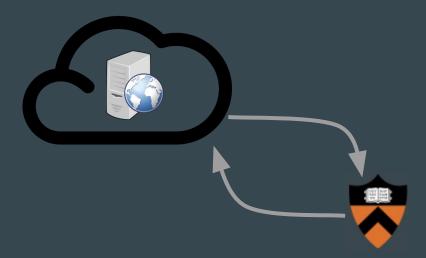
#### Scenario: University working with a small cloud provider

- Small cloud provider offers best price and capabilities
- Performance critical apps need to be reached from campus
- Campus and cloud communicate over the public Internet



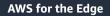
#### Problem: Internet doesn't offer performance guarantees

- Latency can be too large or inconsistent
- Loss can be unacceptably high
- Reliability suffers



#### One Approach: Network performance is offered as a paid service!

Edge computing: performance-critical services placed close to edge networks



Bring the world's most capable and secure cloud to you



**Network-as-a-Service (NaaS)**: on-demand products offering reliable, reserved bandwidth point-to-point links





#### But small networks and underserved regions cannot compete

#### Only hypergiants can afford vast numbers of edge nodes

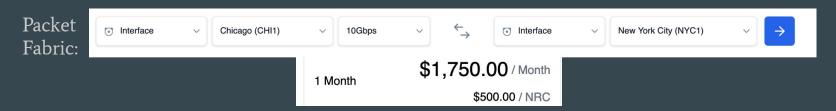
**AWS**:



Google:



#### Only large organizations can afford network-as-a-service



# What would it take for the Internet to serve performance-critical apps?

- We need to know what paths are available (Path Diversity)
- We need to accurately measure performance along those paths (Measurements)
- We need to dynamically route traffic down the best path (Route Control)

### Overcoming the challenges of Internet performance with Tango

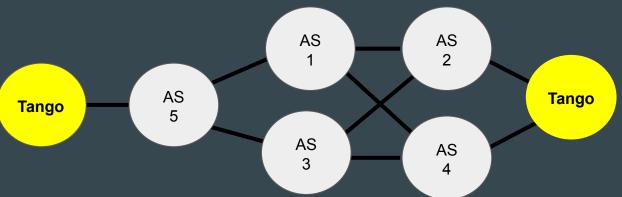


Tango

#### Overcoming the challenges of Internet performance with Tango



Tango



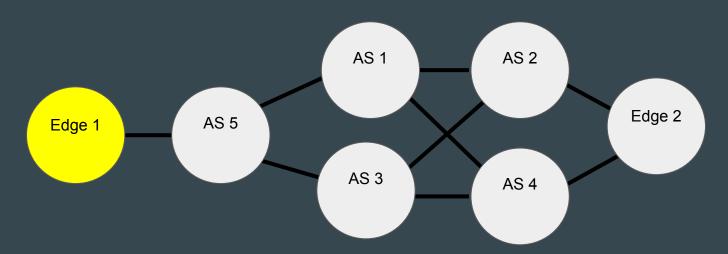
Cooperative and secure edge-to-edge routing with Tango

# What would it take for the Internet to serve performance-critical apps?

- Path Diversity
- **←**

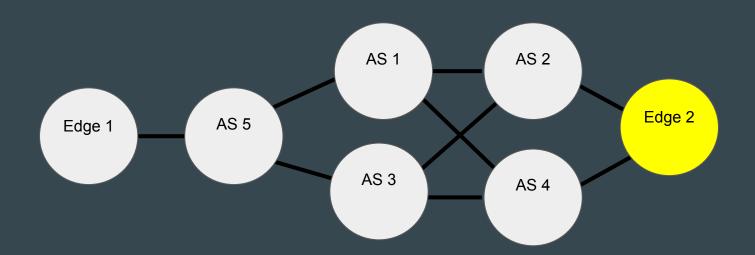
- Measurements
- Route Control

## Why Path Diversity is hard: Traditional Internet Routing (BGP) does not expose all paths



Edge 1 only has a single upstream thus a single path exported by AS 5

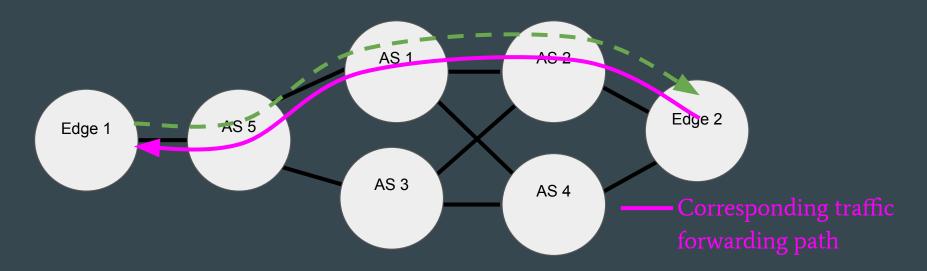
## Why Path Diversity is hard: Traditional Internet Routing (BGP) does not expose all paths



Edge 2 is multi-homed but only influences a single hop, not the whole route

BGP announcement propagation prefix: abcd:1::/48

BGP Pathfinder has no knowledge of the topology:
Begin with the default path

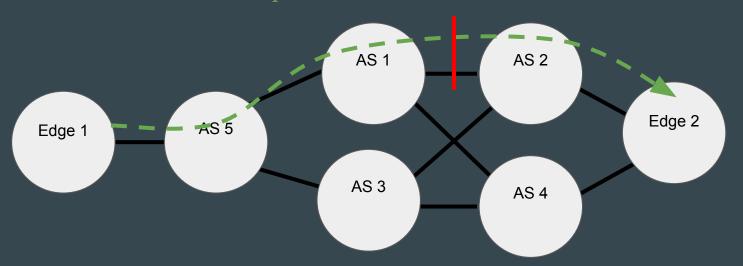


BGP announcement propagation

Suppress the default path

prefix: abcd:1::/48

communities: AS1:No\_Export\_to\_AS2

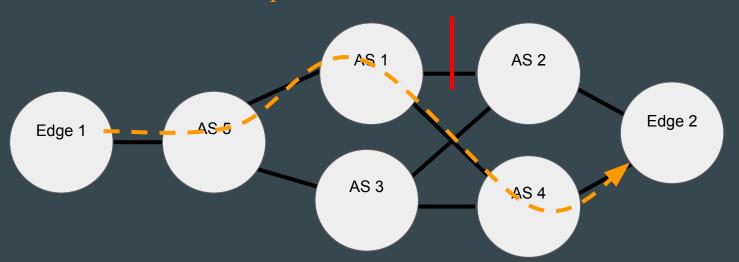


BGP announcement propagation

Find the next path

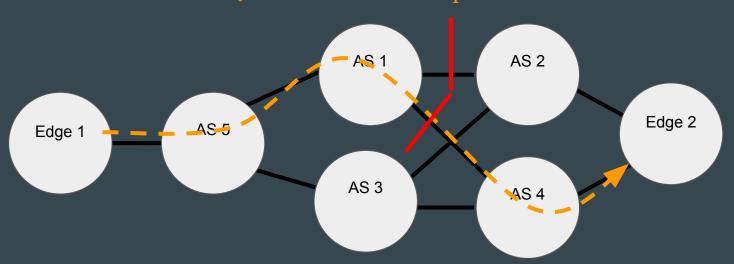
prefix: abcd:1::/48

communities: AS1:No\_Export\_to\_AS2



BGP announcement propagation Suppress the new path prefix: abcd:1::/48

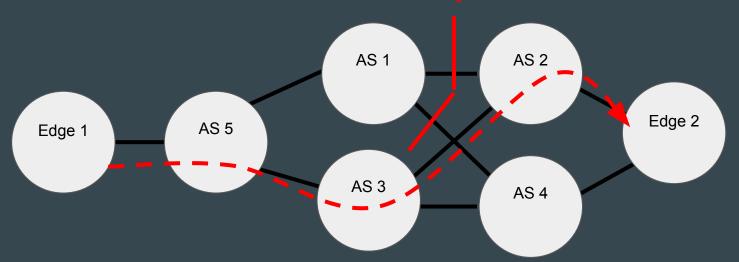
communities: AS1:No\_Export\_to\_AS2, AS1:No\_Export\_to\_AS4



BGP announcement propagation Find the next path

prefix: abcd:1::/48

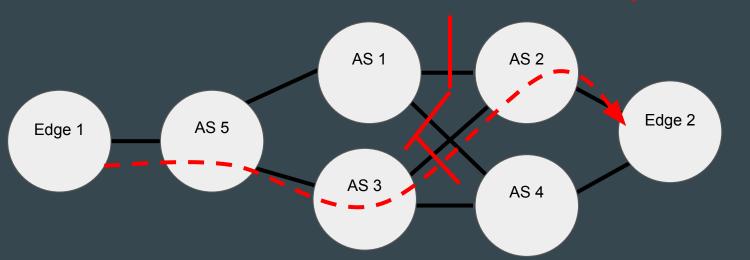
communities: AS1:No\_Export\_to\_AS2, AS1:No\_Export\_to\_AS4



BGP announcement propagation Repeat

prefix: abcd:1::/48

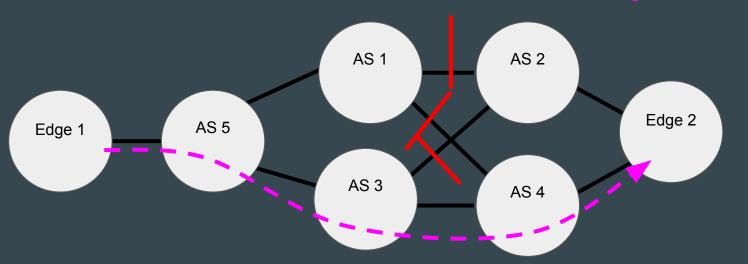
communities: AS1:No\_Export\_to\_AS2, AS1:No\_Export\_to\_AS4, AS3:No\_Export\_to\_AS2



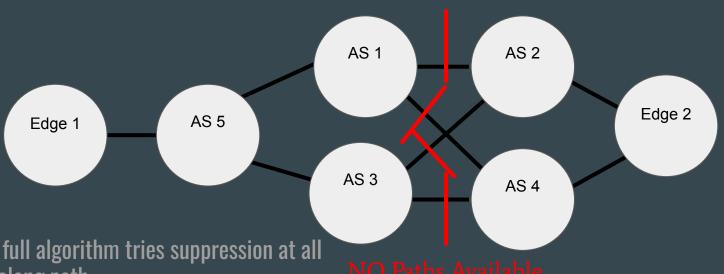
BGP announcement propagation Repeat...

prefix: abcd:1::/48

communities: AS1:No\_Export\_to\_AS2, AS1:No\_Export\_to\_AS4, AS3:No\_Export\_to\_AS2

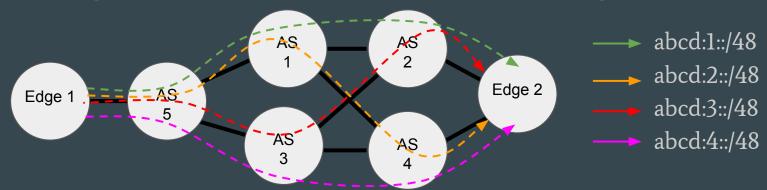


Stop when no paths remain



Note: full algorithm tries suppression at all hops along path

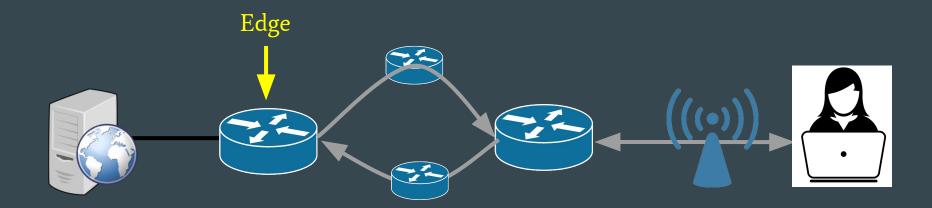
#### Different prefixes are announced along different paths

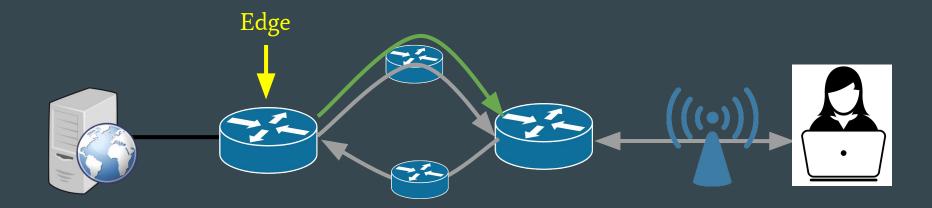


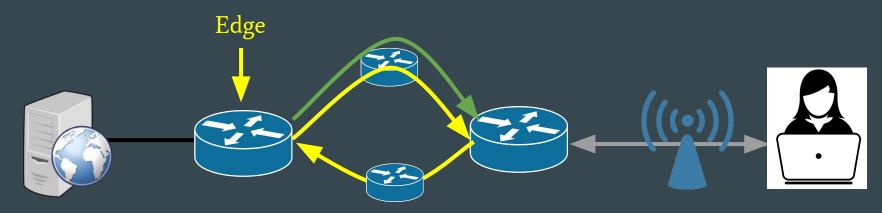
- **Destination** Tango nodes can announce multiple IP prefixes along different paths using BGP communities
- **Source** Tango node can select which path to use by selecting a prefix to reach the destination
- BGP announcements are stable, BGP pathfinder only needs to be rerun periodically 21

# What would it take for the Internet to serve performance-critical apps?

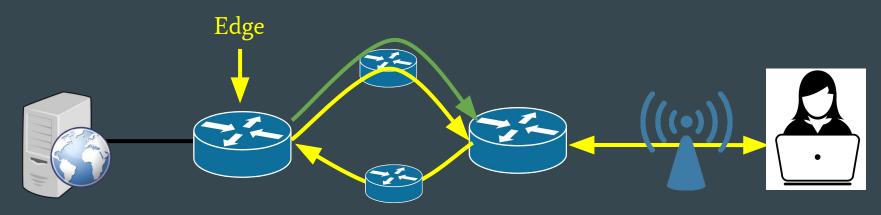
- Path Diversity
- Measurements
- Route Control



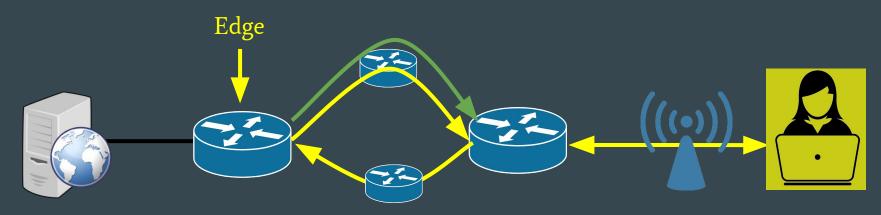




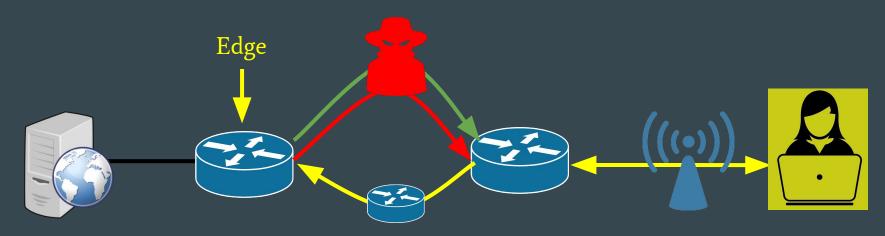
An edge network only sees round-trip-time (RTT) not one-way-delay



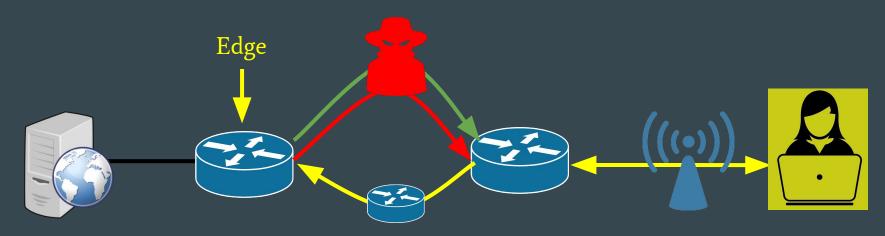
- An edge network only sees round-trip-time (RTT) not one-way-delay
- RTT includes edge network delays



- An edge network only sees round-trip-time (RTT) not one-way-delay
- RTT includes edge network delays
- o RTT monitoring requires understanding end host and L4 protocol behavior



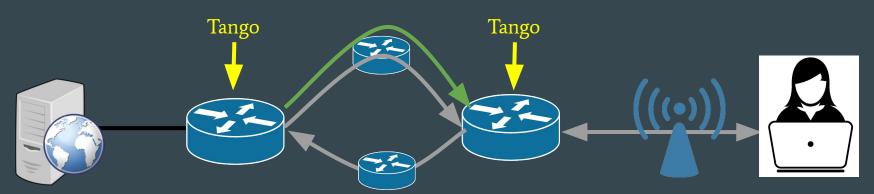
- An edge network only sees round-trip-time (RTT) not one-way-delay
- RTT includes edge network delays
- RTT monitoring requires understanding end host and L4 protocol behavior
- An attacker might try to persuade us that her RTT is lower



- An edge network only sees round-trip-time (RTT) not one-way-delay
- o RTT includes edge network delays
- RTT monitoring requires understanding end host and L4 protocol behavior
- An attacker might try to persuade us that her RTT is lower
- Has to run at line rate

#### Improved Measurements through cooperation

- Custom Tango header added at the edge of the network using programmable switches or eBPF
- Measurements only include one-way, wide-area component
- Measurements do not rely on application behavior



#### **Protecting measurements from adversaries**

An adversary could manipulate measurements to hijack traffic by making her route look more preferable
 How to protect billions of sequence numbers a second?

- Protected packet fields:
  - Sequence numbers:
  - Timestamps:
  - Route control messages:

- protects against adversary hiding lost packets
- protects against adversary manipulating latency
- protects against adversary forging ctrl messages

#### Signing sequence numbers

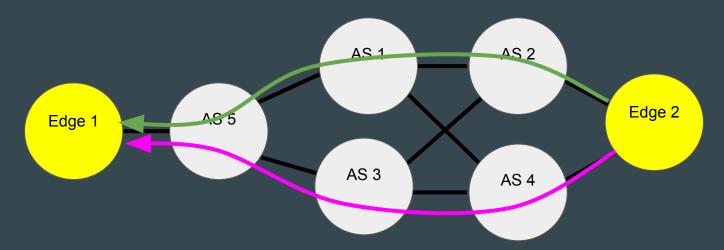
- Each sequence number corresponds to an index in a cipher book shared by the sender and receiver
- Packets contain the sequence number and a single bit "signature" from the corresponding book index to keep up with throughput
- Adversary has a 50/50 chance of guessing a signature, but needs to guess many signatures to meaningfully affect loss
- Likelihood of avoiding detection decreases exponentially with each guess

# What would it take for the Internet to serve performance-critical apps?

- Path Diversity
- Measurements
- Route Control

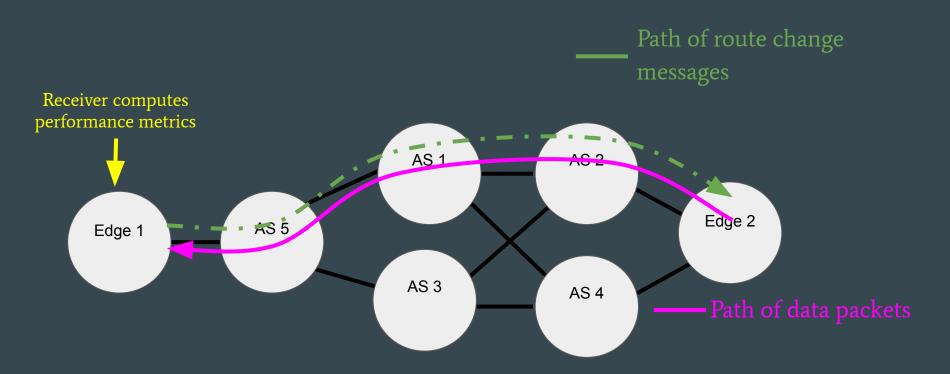
#### Challenges With Route Control

Consider traffic from Edge 2 to Edge 1



- o Edge 1 sees one-way-delay data
- Edge 2 needs to know how to route packets

#### Real-time route control with Tango



#### **Evaluation**

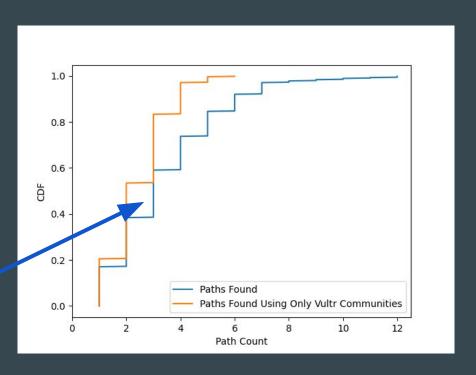
How many paths can we find?



Can we run measurements and crypto at line rate?

#### How many paths can we find?

- Ran BGP pathfinder between 503 globally-distributed nodes from the cloud provider Vultr
- Took ~30min per pair (can be parallelized)
- By default Vultr only exported a single path
- 80% of nodes had additional paths
- BGP Pathfinder can expose 3 paths for the median pair
- Some nodes had as high as 10-12 paths



#### **Evaluation**

- How many paths can we find?
- Can we beat the performance of the default path?



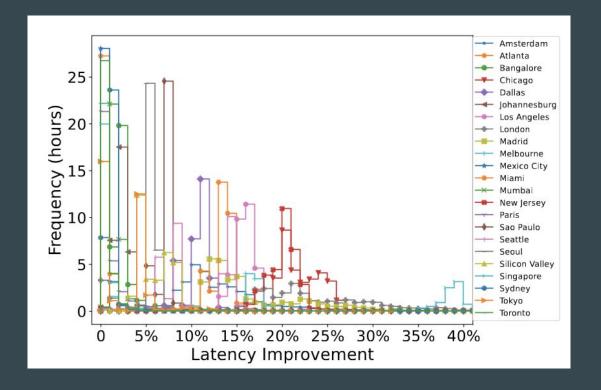
#### Tango performance measurements

- Took measurements from 25 global nodes
- Routed traffic over different paths to two destinations: LA and Stockholm
- Took latency and loss measurement every 10ms



#### **Outperforming the Default Path**

- For many src,dst pairs, the optimal Tango path had 22% lower latency than BGP default path
- In some cases, Tango saw a 39% improvement

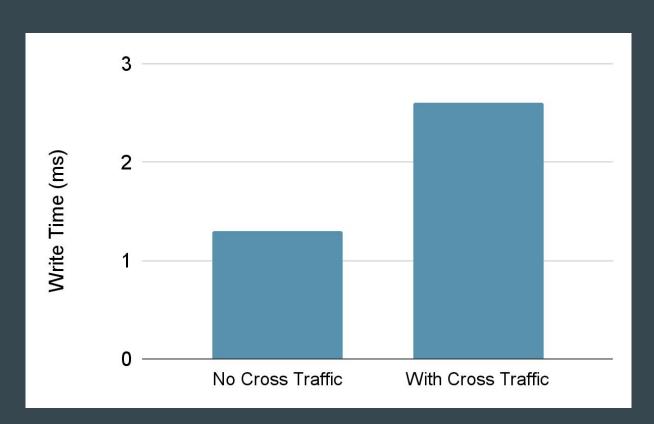


#### **Evaluation**

- How many paths can we find?
- Can we beat the performance of the default path?
- Can we run measurements and crypto at line rate?

#### Generating sequence number signatures at line rate

- Implemented on Tofinol Programmable Switch
- Signature values sent to switch in dataplane
- Switch recirculated packets and wrote signatures to data-plane registers
- Wrote 2<sup>20</sup> signatures in
   2.6ms even with cross traffic
- Keeps up with 100Gbps line rate



#### Conclusion



**Tango** 

3 surprising finds from Tango

- We can find alternative paths through the public Internet
- These paths often have improved performance
- We can run trustworthy telemetry in the data plane

#### **Questions?**

Thank you for your time

Henry Birge-Lee

birgelee@princeton.edu



#### Offering Dynamic Route Control

- Several high-loss and high-latency events plague networks periodically
- Dynamically-switching to better paths can evade these events
- Often other unaffected paths exist

